

Lecture 17 - Nov. 7

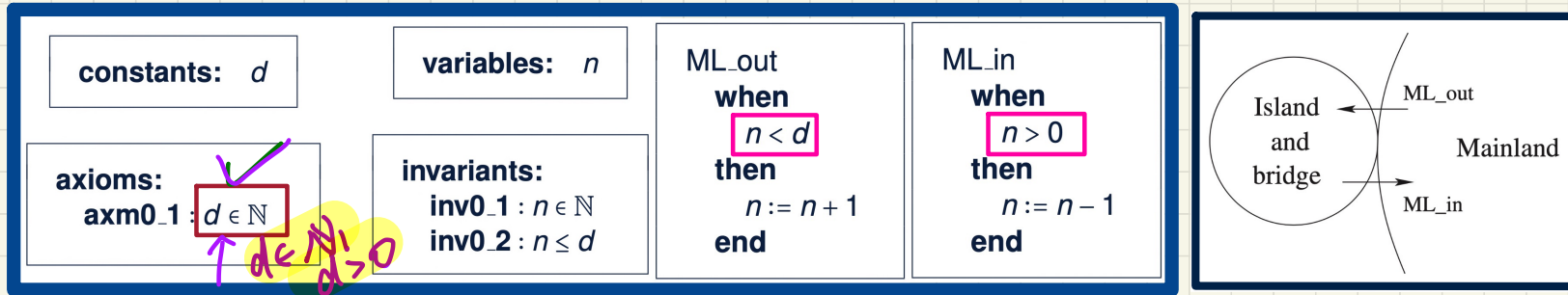
Bridge Controller

Interpreting Unprovable DLF PO
First Refinement: Abstraction, State Space

Announcements/Reminders

- **ProgTest2** results to be released by Monday, Nov 18
- **Lab5** to be released on Friday, Nov 15

Understanding the Failed Proof on DLF



Unprovable Sequent: $\vdash d > 0$

$\vdash d > 0$

Not being able to prove $d > 0$

\hookrightarrow current model may violate $\neg(d > 0)$ is true.

$$\frac{0 < 0}{\text{F}} \vee \frac{0 > 0}{\text{F}} = \text{F}$$

F

deadlocks right after $\neg \text{it}$

Say $d = 0$
 After init : $n = 0$
 $G(\text{ML_out}) \vee G(\text{ML_in})$

- ① $d \leq 0$
- ② $d \in \mathbb{N} \quad (d \geq 0)$

$d = 0$ allowed by current model

Discharging PO of **DLF**: Second Attempt

$$\begin{array}{l} d \in \mathbb{N} \\ n \in \mathbb{N} \\ n \leq d \\ \vdash \\ n < d \vee n > 0 \end{array}$$

$d > 0$ new
condition
added (as a fix).

$$\begin{array}{l} d \in \mathbb{N} \\ n \in \mathbb{N} \\ n < d \vee n = d \\ \vdash \\ n < d \vee n > 0 \end{array}$$

MON

$$\begin{array}{l} d > 0 \\ n < d \vee n = d \\ \vdash \\ n < d \vee n > 0 \end{array}$$

OR_L

$$\begin{array}{l} d > 0 \\ n < d \\ \vdash \\ n < d \vee n > 0 \end{array}$$

OR_R1

$$\begin{array}{l} n < d \\ \vdash \\ n < d \end{array}$$

HYP

$$\begin{array}{l} d > 0 \\ n = d \\ \vdash \\ n < d \vee n > 0 \end{array}$$

EQ_LR, MON

$$\begin{array}{l} d > 0 \\ \vdash \\ d < d \vee d > 0 \end{array}$$

OR_R2

$$\begin{array}{l} d > 0 \\ \vdash \\ d > 0 \end{array}$$

HYP

?

Summary of the Initial Model: Provably Correct

static

constants: d

axioms:

axm0_1 : $d \in \mathbb{N}$

axm0_2 : $d > 0$

fix for the potential deadlock

dynamic

variables: n

invariants:

inv0_1 : $n \in \mathbb{N}$

inv0_2 : $n \leq d$

init

begin

$n := 0$

end

ML_out

when

$n < d$

then

$n := n + 1$

end

added for Inv. Pres.

ML_in

when

$n > 0$

then

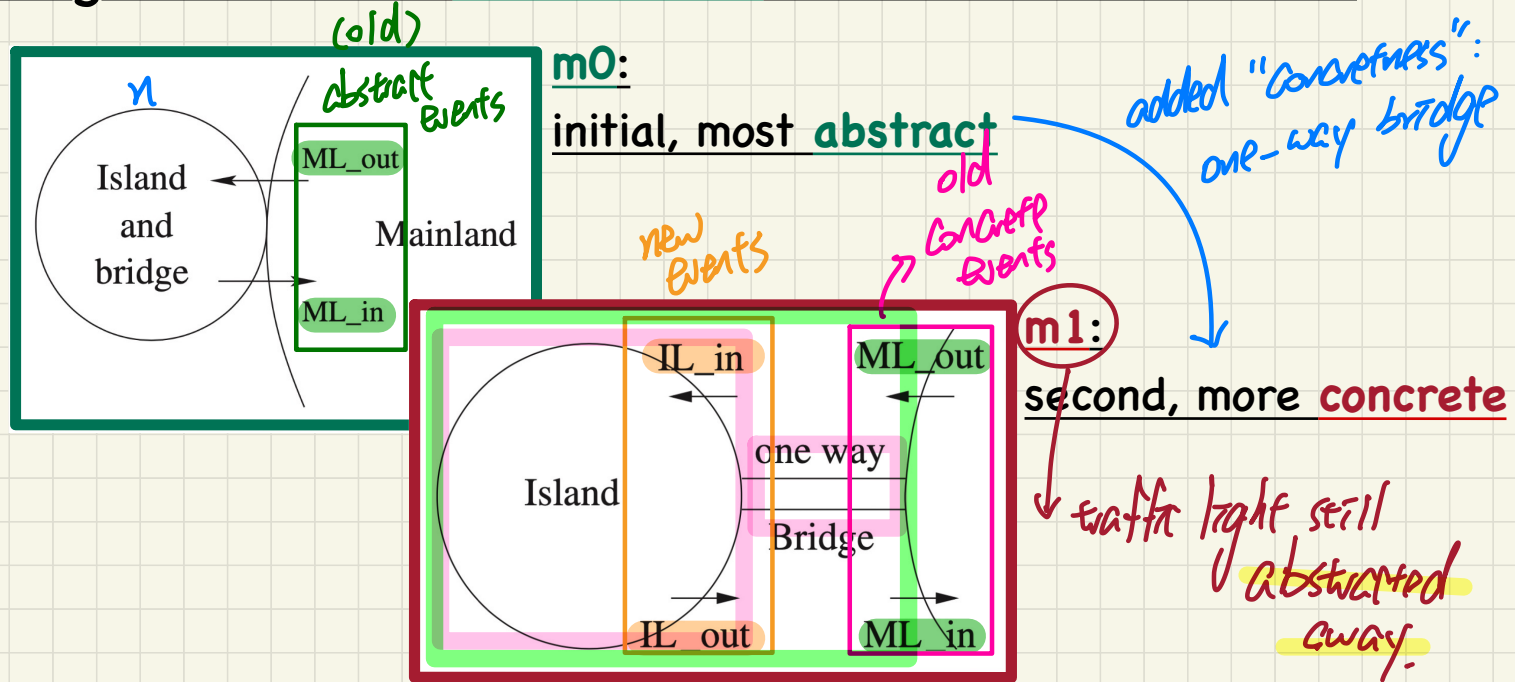
$n := n - 1$

end

Correctness Criteria:

- + Invariant Establishment
- + Invariant Preservation
- + Deadlock Freedom

Bridge Controller: **Abstraction** in the 1st Refinement

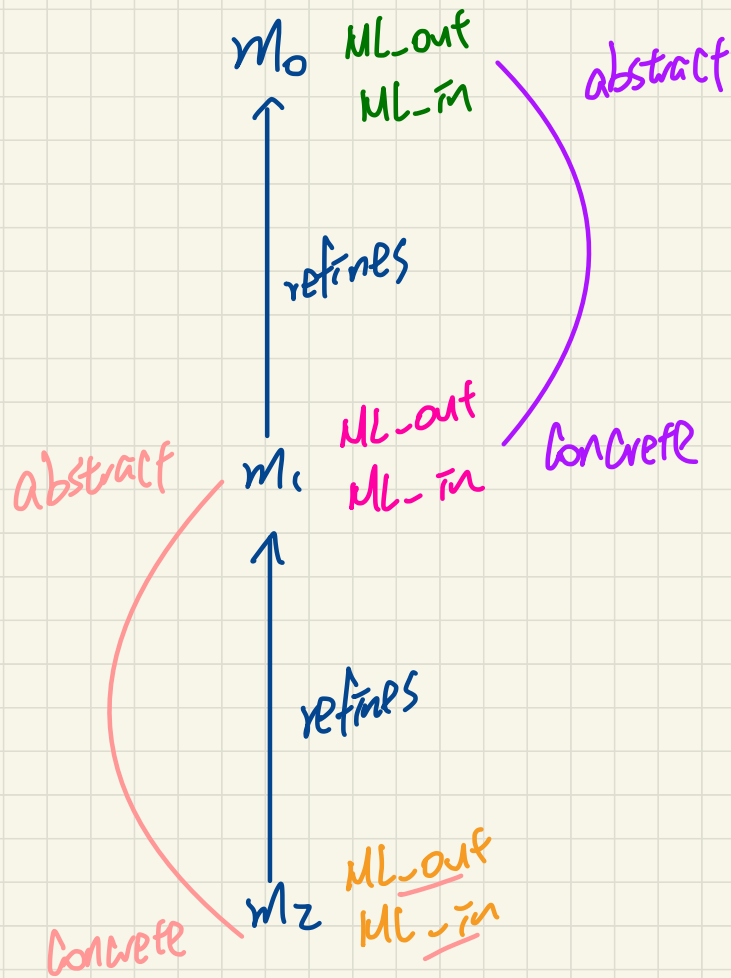


REQ1

The system is controlling cars on a bridge connecting the mainland to an island.

REQ3

The bridge is one-way or the other, not both at the same time.



$$* a=0 \vee c=0 \equiv \neg(a \neq 0 \wedge c \neq 0)$$

Bridge Controller: State Space of the 1st Refinement

REQ1	The system is controlling cars on a bridge connecting the mainland to an island.
REQ3	The bridge is <u>one-way</u> or the other, not both at the same time.

Dynamic Part of Model

variables: a, b, c

invariants:

- inv1.1 : $a \in \mathbb{N}$
- inv1.2 : $b \in \mathbb{N}$
- inv1.3 : $c \in \mathbb{N}$
- inv1.4 : ??
- inv1.5 : ??

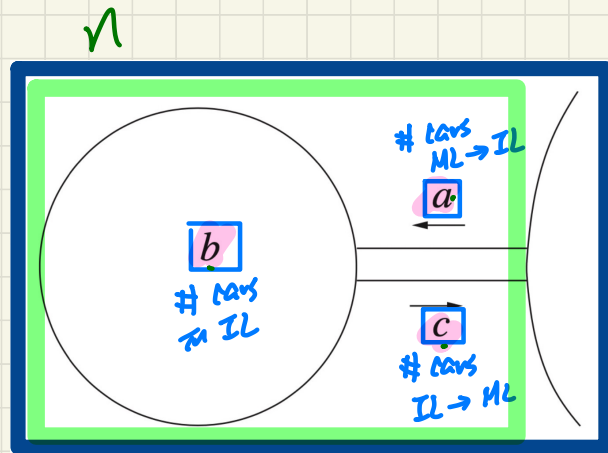
linking/guarding invariant

$$n = a + b + c$$

$$a = 0 \vee c = 0$$

$$a * c = 0$$

at * may be hard for prover



Static Part of Model

constants: d

axioms:

- axm0.1 : $d \in \mathbb{N}$
- axm0.2 : $d > 0$

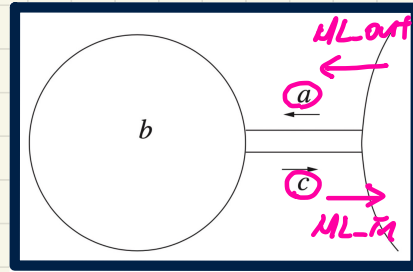
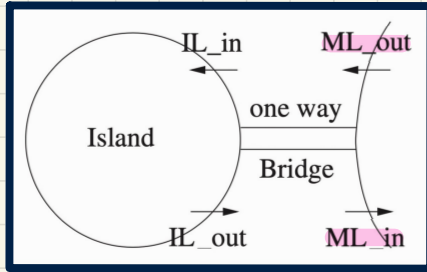
Exercises

inv1.4: linking abstract & concrete states

inv1.5: bridge is one-way

$a \geq b, c$

Bridge Controller: **Guards** of "old" Events 1st Refinement



ML_out: A car exits mainland
(getting on the bridge).

```
ML_out
when
  ??
then
  a := a + 1
end
```

abstract guard from M_0 : $n < d$
(1) From M_0 , abstract guard of ML_out: $n < d$
(2) taking abst. & con. of states: $a + b + c = n$
(3) first concrete guard: $c = 0$

ML_in: A car enters mainland
(getting off the bridge).

```
ML_in
when
  ??
then
  c := c - 1
end
```

constants: d

axioms:

axm0_1 : $d \in \mathbb{N}$
axm0_2 : $d > 0$

(1) $a + b < d$
(2) $a + b < d$
(3) $n < d$
 $a + b + c = 0$

variables: a, b, c

invariants:

inv1_1 : $a \in \mathbb{N}$
inv1_2 : $b \in \mathbb{N}$
inv1_3 : $c \in \mathbb{N}$
inv1_4 : $a + b + c = n$
inv1_5 : $a = 0 \vee c = 0$

Before-After Predicates of Event Actions: 1st Refinement

Events

```
ML_in
  when
     $0 < c$ 
  then
     $c := c - 1$ 
  end
```

```
ML_out
  when
     $a + b < d$ 
     $c = 0$ 
  then
     $a := a + 1$ 
  end
```

- Pre-State
- Post-State
- State Transition

Before-after
predicates

$$a' = a \wedge b' = b \wedge$$
$$c' = c - 1$$
$$a' = a + 1 \wedge b' = b \wedge$$
$$c' = c$$